

ΕΠΛ323 - Θεωρία και Πρακτική Μεταγλωττιστών

Lecture 10a

Type Checking

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Static Checking



- Ensures that certain kids of programming errors will be detected and reported at compile-time:
 - Type checks. An array variable and a function variable are added together.
 - Flow-of-control checks. A break statement in C causes control to leave the smallest enclosing while, for, or switch statement, while the smallest enclosing statement does not exist.
 - Uniqueness checks. Labels in a case statement must be distinct.
 - Name-related checks. In Ada, a loop or block may have a name that appears in the beginning and end of the construct.

Dynamic Checking



- Checks performed by the program at run-time.
- In principle, all checks can be performed at run-time, but this is not efficient.
 - A sound type system allows us to determine statically that these errors cannot occur when the target program runs.
- A language is strongly typed if its compiler can guarantee that a compiled program will execute without errors.
- Not always possible.

```
table: array[0..255] of char;
i: integer
...
x := table[i];
```

Type Systems



- Based on information about the syntactic constructs, the notion of types, and the rules for assigning types to language constructs.
 - "If both operands of the arithmetic operators of addition, subtraction and multiplication are of type integer, then the result is of type integer."
 - "The result of the unary & operator is a pointer of the object referred to by the operand. If the type of the operand is `...', the type of the result is `pointer to...'."

Basic and Constructed Types



- Basic types are atomic types with no internal structure,
 - C: char, int, float, doube, etc.
 - Pascal: boolean, character, integer, real, ranges
 (1..10), enums, etc.
- Constructed types,
 - C: struct, arrays.
 - Pascal: arrays, records, sets.

Type Expressions



- Basic type (boolean, char, integer, and real), and special basic types, type_error, which signals an error during type checking, and void, which denotes the absence of value.
- Type names.
- Type constructors.
 - Arrays.
 - Products.
 - Records.
 - Pointers.
 - Functions.
- Variables holding values that are type expressions.

Simple Type Checker



Examples

```
key: integer;
key mod 1999
array [256] of char
^integer
```

Translation Scheme



```
E \rightarrow D ; E
D \rightarrow D ; D
D \rightarrow id : T \qquad \{ addtype(id.entry, T.type) \}
T \rightarrow char \qquad \{ T.type := char \}
T \rightarrow integer \qquad \{ T.type := integer \}
T \rightarrow ^T_1 \qquad \{ T.type := pointer(T_1.type) \}
T \rightarrow array [num] of T_1 \qquad \{ T.type := array(1..num.val, T_1.type) \}
```

Type Checking of Expressions



```
E \rightarrow literal \{ E.type := char \}
E \rightarrow \text{num}
           { E.type := integer }
E \rightarrow id { E.type := lookup(id.entry) }
E \rightarrow E_1 \mod E_2 \in E.type :=
                   if E_1. type = integer and
                      E_2. type = integer then integer
                      else type error }
E \rightarrow E_1[E_2] { E.type :=
                      if E_2. type = integer and
                          E_1. type = array(s,t) then t
                      else type error }
               { E.type :=
                      if E_1. type = pointer(t) then t
                       else type error }
```

Type Checking of Statements



```
S \rightarrow id : E { S.type :=
                        if id.type = E.type then void
                        else type error }
S \rightarrow \text{if } E \text{ then } S_1 \{ S.type :=
                        if E.type = boolean then <math>S_1.type
                        else type error }
S \rightarrow \text{ while } E \text{ do } S_1 \{ S.type :=
                        if E.type = boolean then <math>S_1.type
                       else type error }
T \rightarrow S_1 ; S_1
                   if S_1. type = void and
                           S_2.type = void then void
                        else type error }
```

Type Checking of Functions



```
E \rightarrow E \ (E) \qquad \{ \text{ $S.type := } \\ \text{ if id.type = $E.type then } void \\ \text{ else } type\_error \ \} \\ T \rightarrow T_1 \ ' \rightarrow ' \ T_2 \qquad \{ \text{ $T.type := } T_1.type \rightarrow T_2.type \ \} \\ E \rightarrow E_1(E_2) \qquad \{ \text{ $E.type := } \\ \text{ if $E_2.type = $s$ and } \\ E_1.type = s \rightarrow t \text{ then } t \\ \text{ else } type\_error \ \}
```